

Quantifying EPR: the resource theory of nonclassicality of common-cause assemblages

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Quantum 7, 926 (2023) and arXiv:2209.10177

Entanglement



- Quantum teleportation
- Quantum Key Distribution

Nonlocality



- Device Independent QKD
- Randomness

Entanglement



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- Quantum Key Distribution

Nonlocality

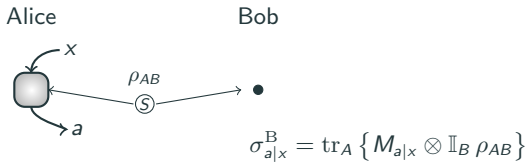


- Device Independent QKD
- Randomness

Einstein-Podolsky-Rosen (EPR) – a.k.a. steering

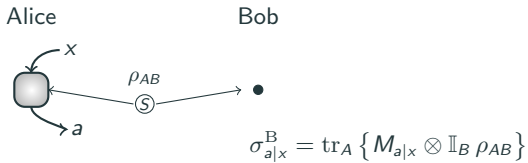


Traditional bipartite setup¹

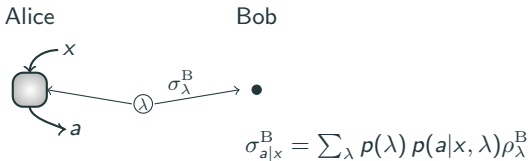


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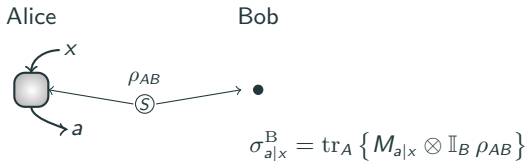


Classical common cause:

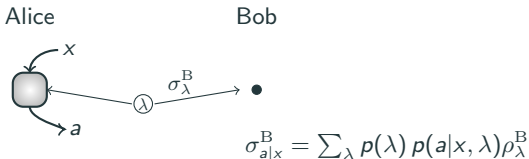


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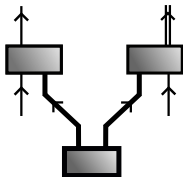


If $\sigma_{a|x}^B$ does not admit of such classical model, then it is EPR non-classical

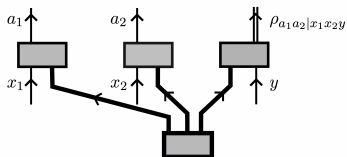
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Generalised EPR scenarios

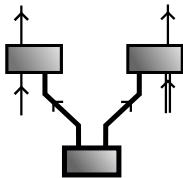
Bob-with-Input: States $\sigma_{a|xy}$



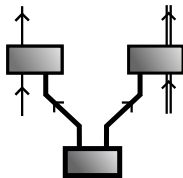
Multipartite Bob-with-Input:
States $\sigma_{a_1 a_2 | x_1 x_2 y}$



Measurement-Device-Independent:
POVM elements $M_{ab|x}$



Channel 'steering':
Quantum channels $\mathcal{I}_{a|x}$



Assemblage: $\{\sigma_{a|xy}\}_{a,x,y}$, $\{\sigma_{a_1 a_2 | x_1 x_2 y}\}_{a_1, a_2, x_1, x_2, y}$, $\{M_{ab|x}\}_{a,b,x}$, $\{\mathcal{I}_{a|x}\}_{a,x}$

Resource Theory I²

Scenario: traditional bipartite and multipartite EPR scenarios.

- Full set of available resources – *Enveloping theory*

Common-cause form, quantum

- Criteria to assert what can be done freely

Key feature: common cause among the wings

free \equiv classical

- Free set of resources and free set of transformations

Processes whose common cause is classical

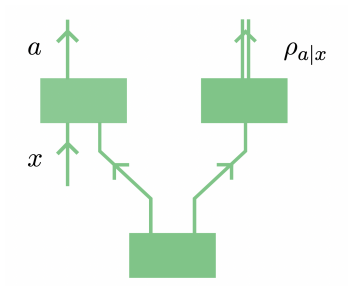
Local Operations and Shared Randomness (LOSR)

²B. Coecke, T. Fritz, R. Spekkens. Information and Computation 250, 59 (2016)

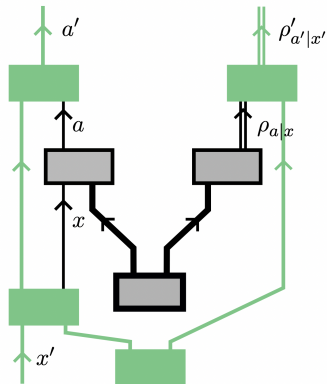
Free resources and transformations

Example: Traditional bipartite EPR scenarios

Free assemblage:



Free Transformation:



Resource Theory II

Scenario: Generalised scenarios – channel EPR scenarios

- Full set of available resources – *Enveloping theory*

Common-cause form, from any theory \longrightarrow non-signalling objects
see Paulo Cavalcanti's talk

- Criteria to assert what can be done freely

Key feature: common cause among the wings

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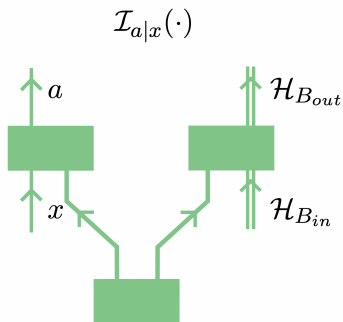
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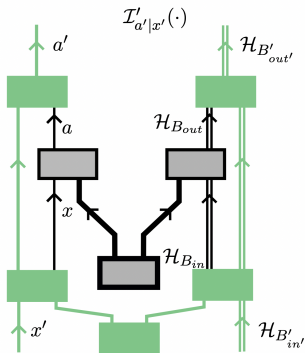
Free resources and transformations

Example: Channel EPR scenarios

Free assemblage:



Free Transformation:



Results

- Result 1: Resource conversion under free operations in all scenarios can be evaluated with a single instance of a semidefinite program.

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Traditional scenarios: (quantum enveloping theory)

SDP 2. $\Sigma_{\mathbb{A}|\mathbb{X}} \xrightarrow{LOSR} \Sigma'_{\mathbb{A}'|\mathbb{X}'}$.

The assemblage $\Sigma_{\mathbb{A}|\mathbb{X}}$ can be converted into the assemblage $\Sigma'_{\mathbb{A}'|\mathbb{X}'}$ under LOSR operations, denoted by $\Sigma_{\mathbb{A}|\mathbb{X}} \xrightarrow{LOSR} \Sigma'_{\mathbb{A}'|\mathbb{X}'}$, if and only if the following SDP is feasible:

$$\begin{aligned} & \text{given } \{ \{ \sigma_{a|x} \}_a \}_x, \{ \{ \sigma'_{a'|x'} \}_{a'} \}_{x'}, \{ D(a'|a, x', \lambda) \}_{\lambda, a', a, x'}, \{ D(x|x', \lambda) \}_{\lambda, x, x'} \\ & \text{find } \{ (W_\lambda)_{BB'} \}_\lambda \\ & \text{s.t. } \begin{cases} W_\lambda \geq 0, \\ \text{tr}_{B'} \{ W_\lambda \} \propto \frac{1}{d} \mathbb{I}_B \quad \forall \lambda, \\ \sum_\lambda \text{tr}_{B'} \{ W_\lambda \} = \frac{1}{d} \mathbb{I}_B, \\ \sigma'_{a'|x'} = \sum_\lambda \sum_{a,x} D(a'|a, x', \lambda) D(x|x', \lambda) d_B \text{tr}_B \{ W_\lambda (\mathbb{I}_{B'} \otimes \sigma_{a|x}^T) \}. \end{cases} \end{aligned} \tag{8}$$

When the conversion is not possible, we denote it by $\Sigma_{\mathbb{A}|\mathbb{X}} \not\xrightarrow{LOSR} \Sigma'_{\mathbb{A}'|\mathbb{X}'}$.

Results

Generalised scenarios: (NS enveloping theory)

SDP 2. *The channel assemblage $\mathbf{I}_{\mathbb{A}|\mathbb{X}}$ can be converted to the channel assemblage $\mathbf{I}'_{\mathbb{A}'|\mathbb{X}'}$ under LOSR operations, denoted by $\mathbf{I}_{\mathbb{A}|\mathbb{X}} \xrightarrow{\text{LOSR}} \mathbf{I}'_{\mathbb{A}'|\mathbb{X}'}$, if and only if the following SDP is feasible:*

$$\begin{aligned}
 & \text{given } \{J_{a|x}\}_{a,x}, \{J'_{a'|x'}\}_{a',x'}, \{D(a'|a, x', \lambda)\}_{\lambda, a', a, x'}, \{D(x|x', \lambda)\}_{\lambda, x, x'} \\
 & \text{find } \{(J_{\xi} \lambda)_{B_{in} B'_{in'}, B_{out} B'_{out'}}\}_{\lambda}, \{(J_F \lambda)_{B_{in} B'_{in'}}\}_{\lambda} \\
 & \text{s.t. } \begin{cases} J_{\xi} \lambda \geq 0 \quad \forall \lambda, \\ \text{tr}_{B'_{out'} B_{in}} \{J_{\xi} \lambda\} \propto \mathbb{I}_{B_{out} B'_{in'}} \quad \forall \lambda, \\ \sum_{\lambda} \text{tr}_{B'_{out'} B_{in}} \{J_{\xi} \lambda\} = \frac{1}{d_{B_{out}} d_{B'_{in'}}} \mathbb{I}_{B_{out} B'_{in'}}, \\ J_F \lambda \geq 0 \quad \forall \lambda, \\ \text{tr}_{B_{in}} \{J_F \lambda\} \propto \mathbb{I}_{B'_{in'}} \quad \forall \lambda, \\ \sum_{\lambda} \text{tr}_{B_{in}} \{J_F \lambda\} = \frac{1}{d_{B'_{in'}}} \mathbb{I}_{B'_{in'}}, \\ \text{tr}_{B'_{out'}} \{J_{\xi} \lambda\} = J_F \lambda \otimes \frac{1}{d_{B_{out}}} \mathbb{I}_{B_{out}} \quad \forall \lambda, \\ J'_{a'|x'} = \sum_{\lambda} \sum_{a,x} D(a'|a, x', \lambda) D(x|x', \lambda) J_{a|x} * J_{\xi} \lambda. \end{cases} \tag{9}
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- Alice's measurements: $x = 0 \rightarrow$ Pauli-Z, $x = 1 \rightarrow$ Pauli-X

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- Family of assemblages: $\{ \Sigma^\theta \mid \theta \in (0, \frac{\pi}{4}] \}$

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Proof: family of monotones $\{ M_\eta \mid \eta \in (0, \frac{\pi}{4}] \}$ such that:

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Hence: $M_\theta \Rightarrow \Sigma^\eta \not\rightarrow \Sigma^\theta$, $M_\eta \Rightarrow \Sigma^\theta \not\rightarrow \Sigma^\eta$

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Bob-with-Input scenario: $\mathcal{H}_B = \text{qubit}$, $\mathbb{X} = \{0, 1, 2\}$, $\mathbb{A} = \{0, 1\}$, $\mathbb{Y} = \{0, 1\}$

$$\Sigma_{\mathbb{A}|\mathbb{X}\mathbb{Y}}^{PTP} = \left\{ \sigma_{a|xy}^{PTP} \right\}_{a \in \mathbb{A}, x \in \mathbb{X}, y \in \mathbb{Y}},$$

$$\text{with } \begin{cases} \sigma_{a|xy}^{PTP} = \xi_y \left\{ \text{tr}_A \left\{ (M_{a|x} \otimes \mathbb{I}_B) |\phi\rangle\langle\phi| \right\} \right\}, \\ M_{a|1} = \frac{\mathbb{I} + (-1)^a \sigma_x}{2}, \quad M_{a|2} = \frac{\mathbb{I} + (-1)^a \sigma_y}{2}, \quad M_{a|3} = \frac{\mathbb{I} + (-1)^a \sigma_z}{2}, \end{cases}$$

$$\Sigma_{\mathbb{A}|\mathbb{X}\mathbb{Y}}^{PR} = \left\{ \sigma_{a|xy}^{PR} \right\}_{a \in \mathbb{A}, x \in \mathbb{X}, y \in \mathbb{Y}},$$

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$$\Sigma^{PTP} \not\rightarrow \Sigma^{PR} \quad \text{and} \quad \Sigma^{PTP} \not\leftarrow \Sigma^{PR}$$

Comparison with other work

Traditional EPR scenario – multipartite:

Free assemblage: $\sigma_{a_1 a_2 | x_1 x_2} = \sum_{\lambda} p(\lambda) \rho(a_1 a_2 | x_1 x_2 \lambda) \rho_{\lambda}$

$\rho(a_1 a_2 | x_1 x_2 \lambda)$ can be signalling!

Free operations: wirings and stochastic 1-way LOCC

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free operation on free resource gives non-free assemblage

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Need for a principled approach to ‘freeness’

Conclusions and outlook

- Traditional and generalised EPR scenarios
- Causal perspective on their 'object of interest' (assemblages)
- Causally-underpinned resource theory:
 - free = classical common cause (LOSR)
- (i) Traditional scenarios + quantum enveloping theory
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- Resource conversion: tested by a single SDP
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