

The Algebra for Stabilizer Codes

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Outline

The ZX-calculus

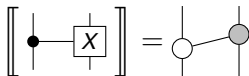
Graphical linear/affine algebra and the ZX-calculus

Graphical affine Lagrangian algebra and stabilizer circuits
[Comfort and Kissinger, 2022]

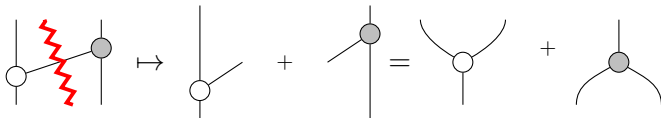
Graphical affine coisotropic algebra and stabilizer codes
[Comfort, 2023]

The ZX-calculus “splits the atom”

The ZX-calculus decomposes quantum circuits into smaller components. Consider the controlled-X gate:



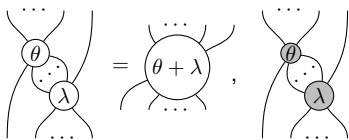
“ **Splitting the atom** ”



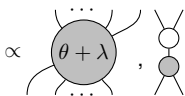
These components are no longer circuits, but they are useful.

Spider fusion and Hopf law

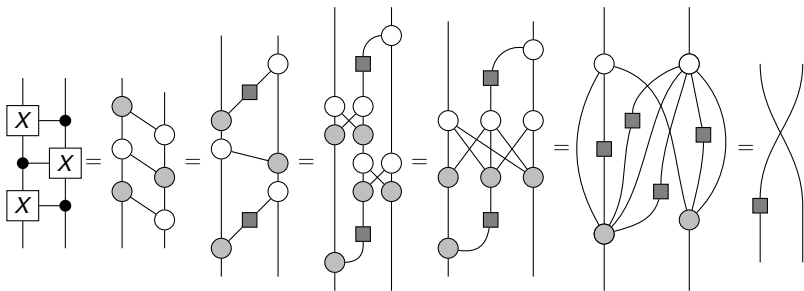
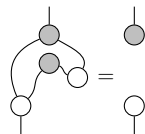
Spider fusion:



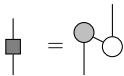
Bialgebra:



Hopf law:



Where



is the antipode

Interpreting the ZX-calculus

The d dimensional qudit (pure) **ZX-calculus** are a family of graphical languages for $d^n \times d^m$ dimensional complex matrices.

There are two families of generators, Z and X spiders, decorated by phases $\vec{\theta} = (0, \theta_1, \theta_2, \dots, \theta_{d-1}) \in [0, 2\pi)^d$:

$$\left[\left[\begin{array}{c} m \\ \vdots \\ \theta \\ \vdots \\ n \end{array} \right] \right] = \frac{1}{\sqrt{d}} \sum_{j=0}^{d-1} e^{i\theta_j} |j, \dots, j\rangle \langle j, \dots, j|$$

$$\left[\left[\begin{array}{c} m \\ \vdots \\ \theta \\ \vdots \\ n \end{array} \right] \right] = \frac{1}{\sqrt{d}} \mathcal{F}^\dagger \sum_{j=0}^{d-1} e^{i\theta_j} |j, \dots, j\rangle \langle j, \dots, j| \mathcal{F}$$

Terminology:

qudit \rightarrow **qubit** when $d = 2$.

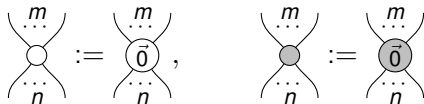
qudit \rightarrow **quopit** when is an d odd prime.

Fragments of the ZX-calculus

Fragments of the ZX-calculus have restricted phases:

Example

The **qudit phase-free fragment** ZX-calculus has trivial angles:



The **qubit stabilizer fragment** has Z/X angles in:

$$\{(0, 0), (0, \pi/2), (0, \pi), (0, 2\pi/3)\} \subseteq [0, 2\pi)^2$$

The **quopit stabilizer fragment** has Z/X angles in:

$$\left\{ \prod_{j=0}^{p-1} (n \cdot j + m \cdot j^2) \pi / p \mid \forall n, m \in \mathbb{F}_p \cong \mathbb{Z}/p\mathbb{Z} \right\} \subseteq [0, 2\pi)^p$$

Linear relations

Definition

Given a field k there is a prop LinRel_k of **linear relations**, whose maps $n \rightarrow m$ are linear subspaces of $k^n \oplus k^m$ under relational composition: For $S \subseteq k^n \oplus k^m$ and $R \subseteq k^m \oplus k^\ell$

$$S; R := \{(x, z) \in k^n \oplus k^\ell \mid \exists y \in k^m : (x, y) \in S \wedge (y, z) \in R\}$$

Lemma

LinRel_k is generated by two spiders, scalars (+ equations):

$$\left[\begin{array}{c} \cdots \\ \text{---} \\ \bigcirc \\ \text{---} \\ \cdots \end{array} \right] = \left\{ \left(\left(\begin{array}{c} a \\ \vdots \\ a \end{array} \right), \left(\begin{array}{c} a \\ \vdots \\ a \end{array} \right) \right) \mid \forall a \in k \right\}$$

$$\left[\begin{array}{c} \cdots \\ \text{---} \\ \bullet \\ \text{---} \\ \cdots \end{array} \right] = \left\{ \left(\left(\begin{array}{c} a_1 \\ \vdots \\ a_n \end{array} \right), \left(\begin{array}{c} b_1 \\ \vdots \\ b_m \end{array} \right) \right) \mid \forall a_i, b_j \in k : \sum_{i=1}^n a_i = \sum_{j=1}^m b_j \right\}$$

$$\left[\begin{array}{c} | \\ \text{---} \\ \text{c} \\ \text{---} \\ | \end{array} \right] = \{(a, a \cdot c) \mid \forall a \in k\}$$

Lemma ([Zanasi, 2018])

$\text{LinRel}_{\mathbb{F}_p}$ is isomorphic to the p -dimensional qupit phase-free ZX-calculus modulo scalars.

Proof.

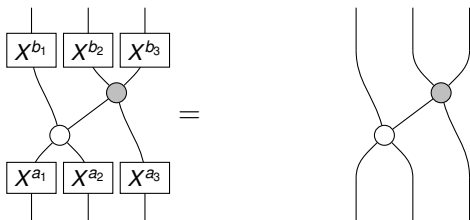
A phase free ZX-diagram D is identified with its X stabilizers:

$$\left[\left[\begin{array}{c} \dots \\ \text{---} \\ D \\ \text{---} \\ \dots \\ n \end{array} \right] \right]_X := \left\{ \left(\left(\begin{array}{c} a_1 \\ \vdots \\ a_n \end{array} \right), \left(\begin{array}{c} b_1 \\ \vdots \\ b_m \end{array} \right) \right) \in \mathbb{F}_p^n \oplus \mathbb{F}_p^m \mid \begin{array}{c} \text{---} \dots \text{---} \\ X^{b_1} \ X^{b_m} \ X^{a_1} \ X^{a_n} \\ \text{---} \dots \text{---} \\ \text{---} \end{array} \right\} = \left[\begin{array}{c} \dots \\ \text{---} \\ D \\ \text{---} \\ \dots \end{array} \right]_X$$

□

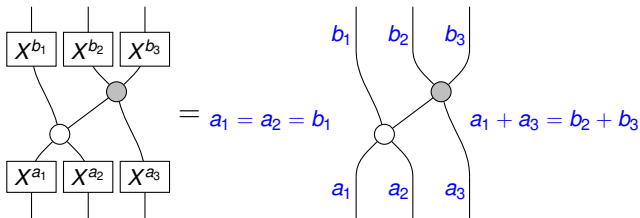
Example: part i

Find the $a_1, a_2, a_3, b_1, b_2, b_3$ which satisfy:



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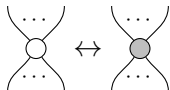


These equations determine a linear subspace of $\mathbb{F}_p^3 \oplus \mathbb{F}_p^3$:

$$\left\{ \left(\left(\begin{pmatrix} a_1 \\ a_2 \\ a_3 \end{pmatrix}, \begin{pmatrix} b_1 \\ b_2 \\ b_3 \end{pmatrix} \right) \mid a_1 = a_2 = b_1 \wedge a_1 + a_3 = b_2 + b_3 \right\}$$

What about the Z stabilizers?

The **Fourier transform** of a ZX-diagram is the colour swapping:



In the phase-free ZX-calculus, this corresponds to the **orthogonal complement** of linear subspaces:

$$(_)^\perp : \text{LinRel}_{\mathbb{F}_p} \rightarrow \text{LinRel}_{\mathbb{F}_p};$$

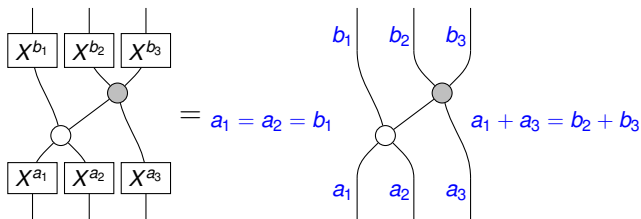
$$(\mathcal{S} \subseteq \mathbb{F}_p^n) \mapsto (\{v \in \mathbb{F}_p^n \mid \forall w \in \mathcal{S}, v^T w = 0\} \subseteq \mathbb{F}_p^n)$$

So we can also identify phase-free ZX-diagrams D with their Z stabilizers $\llbracket D \rrbracket_Z$:

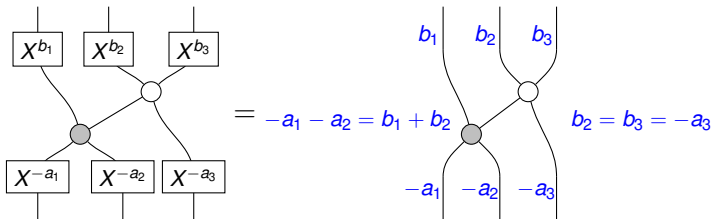
$$\llbracket D \rrbracket_Z = \llbracket D \rrbracket_X^\perp$$

Example: part ii

Recall how we calculated the X stabilizers:



For Z stabilizers, find X stabilizers of the Fourier transform:



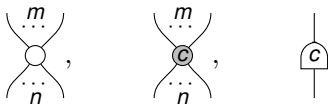
Affine relations

Definition ([Bonchi et al., 2019])

There is a prop AffRel_k of **affine relations**, but now the morphisms are (possibly empty) *affine* subspaces of $k^n \oplus k^m$.

Lemma

AffRel_k is presented by decorating the grey spiders of LinRel_k with all elements $c \in k$:



with the interpretation:

$$\left[\left[\begin{array}{c} \dots \\ \text{c} \\ \dots \end{array} \right] \right] = \left\{ \left(\left(\begin{array}{c} a_1 \\ \vdots \\ a_n \end{array} \right), \left(\begin{array}{c} b_1 \\ \vdots \\ b_m \end{array} \right) \right) \mid \forall a_i, b_j \in k : c + \sum a_i = \sum b_j \right\}$$

Modulo even more equations.

X-phase fragment of ZX and affine relations

The X gate is a phase:

$$\boxed{X^a} = \textcircled{\vec{\theta}}, \quad \text{where } \vec{\theta} = (0a\pi/p, 1a\pi/p, 2a\pi/p, \dots, (p-1)a\pi/p)$$

So consider the **phase-free+X fragment** of the ZX-calculus.

Theorem

$\text{AffRel}_{\mathbb{F}_p}$ is isomorphic to the p -dimensional qudit phase free+X ZX-calculus gate modulo scalars.

Proof.

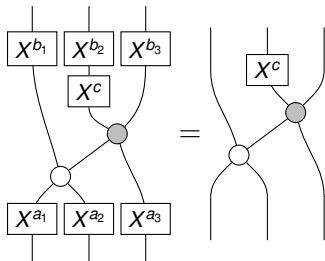
The affine shift and X -gate have the same effect on X stabilizers.

$$X^a = \sum_{x=0}^{p-1} |x+a\rangle\langle x|, \quad \left[\textcircled{a} \right] = \{(x, x+a)\}$$



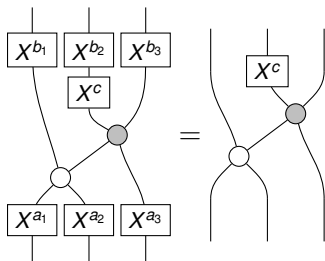
Example: part iii

What are the $a_1, a_2, a_3, b_1, b_2, b_3 \in \mathbb{F}_p$ which satisfy:

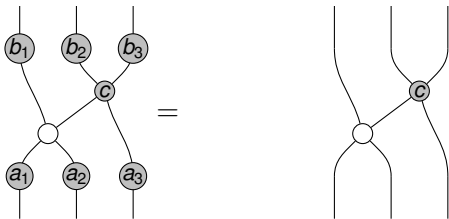


Example: part iii

What are the $a_1, a_2, a_3, b_1, b_2, b_3 \in \mathbb{F}_p$ which satisfy:

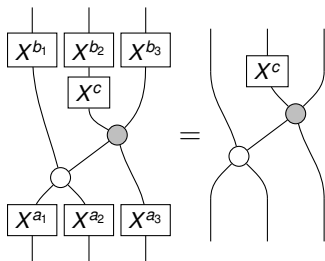


We can translate this into affine relations:

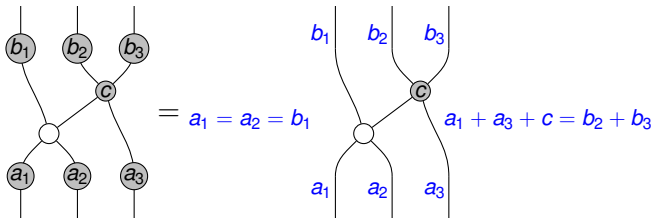


Example: part iii

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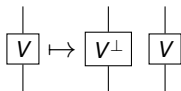
We can translate this into affine relations:



Adding Z and X gates to the picture

We can combine the Z and X -stabilizers in the same picture.

The phase-free circuits are doubled linear relations:



And add *both* the Z and X gates at once:

$$\left[\begin{array}{c} | \\ \textcircled{a} \\ | \end{array} \right]_{\text{Mat}_{\mathbb{C}}/\sim} = Z^a \qquad \left[\begin{array}{c} | \\ | \\ \textcircled{a} \\ | \end{array} \right]_{\text{Mat}_{\mathbb{C}}/\sim} = X^a$$

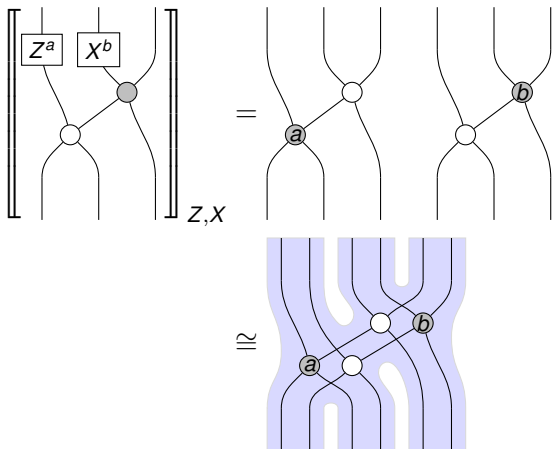
The Z/X gates shift the Z/X stabilizers:

$$\left[\begin{array}{c} | \\ \textcircled{a} \\ | \end{array} \right]_{Z,X} = \left\{ \left(\begin{pmatrix} z \\ x \end{pmatrix}, \begin{pmatrix} z+a \\ x \end{pmatrix} \right) \right\}, \quad \left[\begin{array}{c} | \\ | \\ \textcircled{a} \\ | \end{array} \right]_{Z,X} = \left\{ \left(\begin{pmatrix} z \\ x \end{pmatrix}, \begin{pmatrix} z \\ x+a \end{pmatrix} \right) \right\}$$

This captures the **phase-free+Z+X** fragment of the ZX calculus

Example: part iv

Consider the interpretation of this ZX diagram:



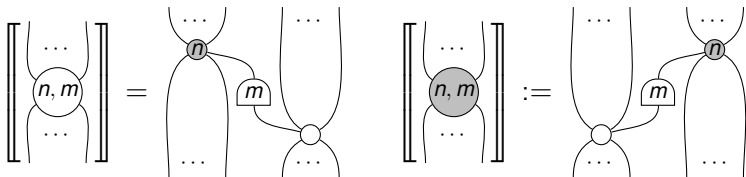
The Z and X stabilizers don't interact with each other for phase-free $+Z+X$ circuits...

Getting all stabilizer circuits

We can add the phase-shift gates to the picture:

Theorem ([Comfort and Kissinger, 2022])

Quopit stabilizer circuits are generated by the affine relations:



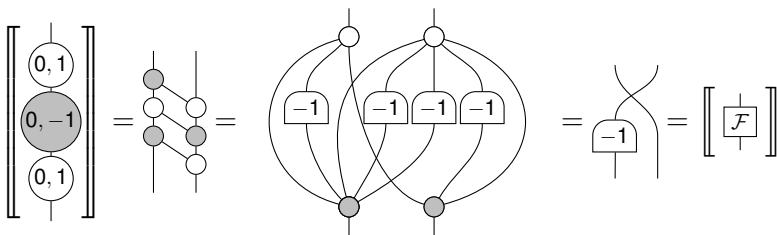
The white spider has the following interpretation in $\text{Mat}_{\mathbb{C}}$:

$$\left[\begin{array}{c} \dots \\ \text{white spider } (n, m) \\ \dots \end{array} \right] = \frac{1}{\sqrt{p}} \sum_{j=0}^{p-1} e^{\pi \cdot i(n \cdot j + m \cdot j^2)/p} |j, \dots, j\rangle \langle j, \dots, j|$$

The interpretation of the grey spider is analogous.

“Splitting the atom” $\times 2$

The Fourier transform is derived by Euler composition:



And the Euler decomposition is derived from the Hopf law!

Symplectic algebra and Weyl operators

Definition

Given $(z_1, \dots, z_n, x_1, \dots, x_n) \in \mathbb{F}_p^{2n}$, there is a **Weyl operator**:

$$\mathcal{W}(z, x) := \bigotimes_{j=1}^n z_{(j)}^{z_j} x_{(j)}^{x_j} : \mathbb{C}^n \rightarrow \mathbb{C}^n$$

The **symplectic form** $\omega : \mathbb{F}_p^{2n} \oplus \mathbb{F}_p^{2n} \rightarrow \mathbb{F}_p$ takes

$$\left(\begin{pmatrix} z \\ x \end{pmatrix}, \begin{pmatrix} z' \\ x' \end{pmatrix} \right) \mapsto z^T x' - x'^T z$$

This form captures the commutation of Weyl operators:

$$\mathcal{W}(z, x)\mathcal{W}(z', x') = e^{i \cdot \pi \cdot \omega((z, x), (z', x')) / p} \mathcal{W}(z', x')\mathcal{W}(z, x)$$

Stabilizer circuits are affine Lagrangian relations

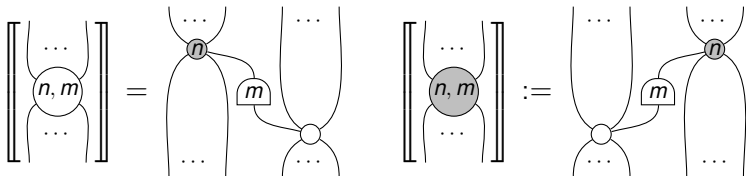
Definition

Lagrangian subspaces are linear subspaces $V \subseteq \mathbb{F}_p^{2n}$ st:

$$V = V^\omega := \{v \in \mathbb{F}_p^{2n} \mid \forall w \in V, \omega(v, w) = 0\}$$

These are linear subspaces of \mathbb{F}_p^{2n} where all elements commute wrt the symplectic form.

The two spiders we mentioned actually generate the prop $\text{AffLagRel}_{\mathbb{F}_p}$ of **affine Lagrangian relations** over \mathbb{F}_p :

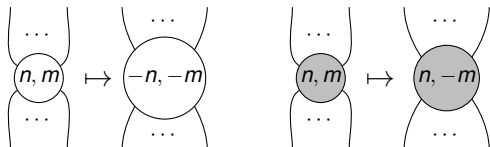


The view of stabilizer *states* in terms of affine Lagrangian subspaces acted on by reversible transformations with measurement statistics was shown in [Calderbank et al., 1998, Gross, 2006, Catani and Browne, 2017, De Beaudrap, 2013]

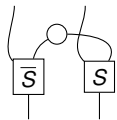
Mixed stabilizers

The complex conjugation is given by

$$\overline{(_)} : \text{AffLagRel}_{\mathbb{F}_p} \rightarrow \text{AffLagRel}_{\mathbb{F}_p};$$



Mixed stabilizer circuits are constructed by doubling stabilizer circuits S :



The quantum discard is the cap: $\left[\begin{array}{c} \equiv \\ \top \\ \equiv \\ \equiv \end{array} \right] = \text{cap}$

Formally this is taking $\text{CPM}(\text{AffLagRel}_{\mathbb{F}_p}, \overline{(_)}) \cong \text{CPM}(\text{Stab}_p, \overline{(_)}) / \sim$

Stabilizer codes

Definition

There is a category of **affine coisotropic relations**

$\text{AffColsotRel}_{\mathbb{F}_p}$ whose morphisms $n \rightarrow m$ are affine subspaces

$$L + a \subseteq \mathbb{F}_p^{2(n+m)} \text{ such that } L^\omega \subseteq L$$

(relaxed from affine Lagrangian subspaces where $L^\omega = L$)

Proposition ([Comfort, 2023])

$\text{AffColsotRel}_{\mathbb{F}_p}$ is presented by adding the discard relation to $\text{AffLagRel}_{\mathbb{F}_p}$:

$$\llbracket \begin{array}{c} \circ \\ | \\ \circ \end{array} \begin{array}{c} \circ \\ | \\ \circ \end{array} \rrbracket = \left\{ \left(\begin{pmatrix} z \\ x \end{pmatrix}, * \right) \mid \forall z, x \in \mathbb{F}_p \right\}$$

Theorem ([Comfort, 2023])

$$\text{AffColsotRel}_{\mathbb{F}_p} \cong \text{CPM}(\text{AffLagRel}_{\mathbb{F}_p}) \cong \text{CPM}(\text{Stab}_p) / \sim$$

The quantum discard is the discard relation!

(reminiscent of [Carette et al., 2021])

Measurement

The Z and X projectors look as follows:

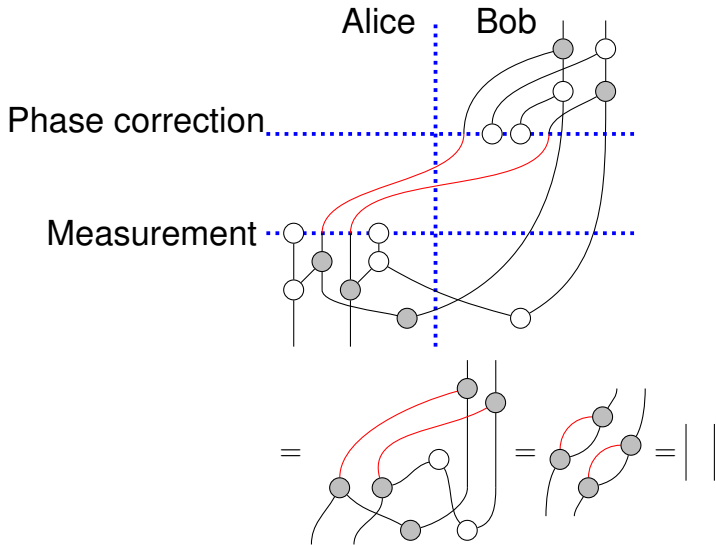
$$\rho_Z := \begin{array}{c} \text{---} \\ | \quad \circ \\ | \quad / \\ | \quad \circ \\ | \quad | \end{array} \quad \begin{array}{c} \text{---} \\ | \quad \circ \\ | \quad \backslash \\ | \quad \circ \\ | \quad | \end{array} = \left| \begin{array}{c} \circ \\ \circ \end{array} \right|, \quad \rho_X := \begin{array}{c} \text{---} \\ | \quad \circ \\ | \quad \backslash \\ | \quad \circ \\ | \quad | \end{array} \quad \begin{array}{c} \text{---} \\ | \quad \circ \\ | \quad / \\ | \quad \circ \\ | \quad | \end{array} = \left| \begin{array}{c} \circ \\ \circ \end{array} \right|$$

By splitting the X projector, following [Selinger, 2008], we get state preparation and measurement relations:

$$\left[\begin{array}{c} \circ \\ | \end{array} \right] = \left\{ \left(x, \begin{pmatrix} z \\ x \end{pmatrix} \right) \mid \forall z, x \in \mathbb{F}_p \right\}$$
$$\left[\begin{array}{c} \circ \\ | \end{array} \right] = \left\{ \left(\begin{pmatrix} z \\ x \end{pmatrix}, x \right) \mid \forall z, x \in \mathbb{F}_p \right\}$$

Possible outcomes of Pauli measurements of stabilizer states are equally likely.

Quantum teleportation by spider fusion



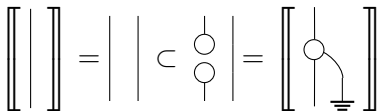
(this also works for qubits)

Maps between stabilizer codes

This is a 2-category with respect to subspace inclusion:

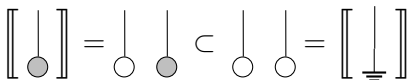
Example

All stabilizers for the identity channel are stabilizers for the Z/X -projectors:



Example

All stabilizers for $|0\rangle$ are stabilizers for the maximally mixed state:

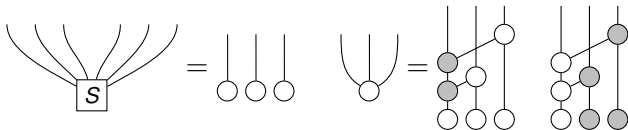


Error correction example: repetition code

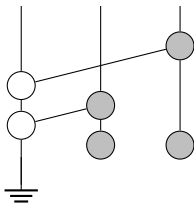
Consider the Linear subspace:

$$S = \{((z_1, z_2, z_3), (x_1, x_2, x_3)) \mid x_1 = x_2 = x_3\} \subseteq \mathbb{F}_2^{2(3)}$$

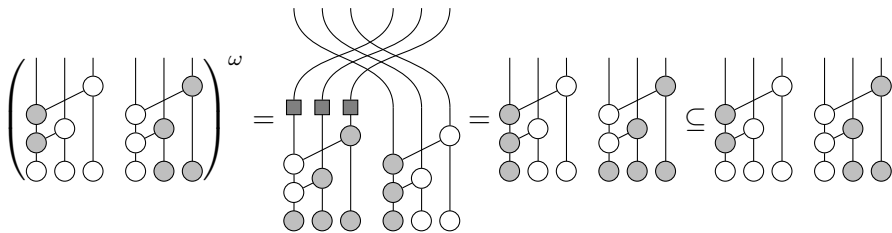
Which can be written in the form of a circuit:



In the undoubled picture this looks like:

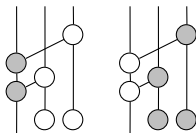


Recall S is coisotropic if $S^\omega \subseteq S$:



Moreover $S \subseteq \mathbb{F}_2^{2(3)}$ has dimension $3 + 1 = 4$

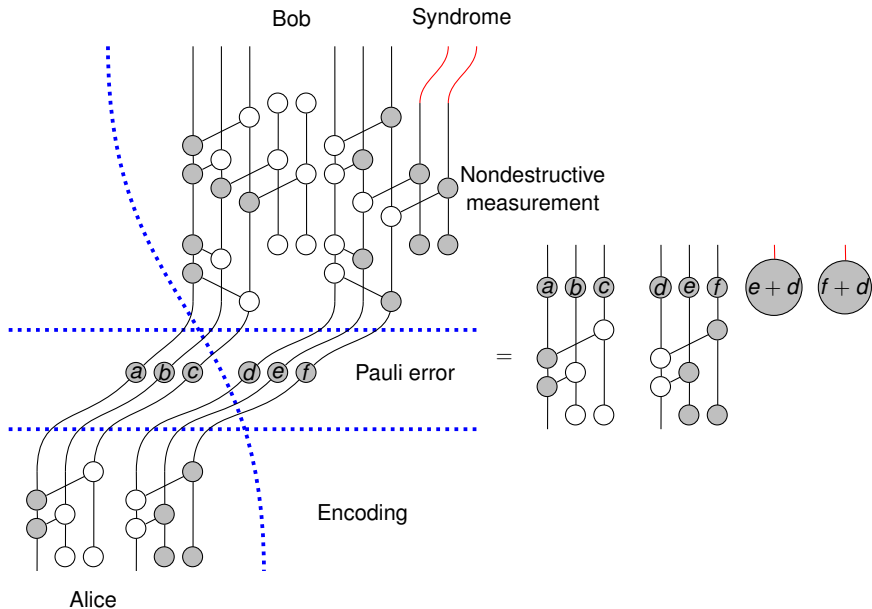
Chopping off the maximally mixed state gives us an encoding map $\mathbb{F}_2^{2(1)} \rightarrow \mathbb{F}_2^{2(3)}$:



That is, this is a qubit $[3, 1]$ -stabilizer code.

Encodes 1 logical qubit into 3 physical qubits.

Suppose there is a Pauli error $W((a, b, c), (d, e, f))$, then we have the following error detection circuit:



Suppose we want to correct for single pauli X errors, then we find that:

$$X \text{ error } (d, e, f) = (1, 0, 0) \text{ syndrome } (e + d, f + d) = (1, 1)$$

$$X \text{ error } (d, e, f) = (0, 1, 0) \text{ syndrome } (e + d, f + d) = (1, 0)$$

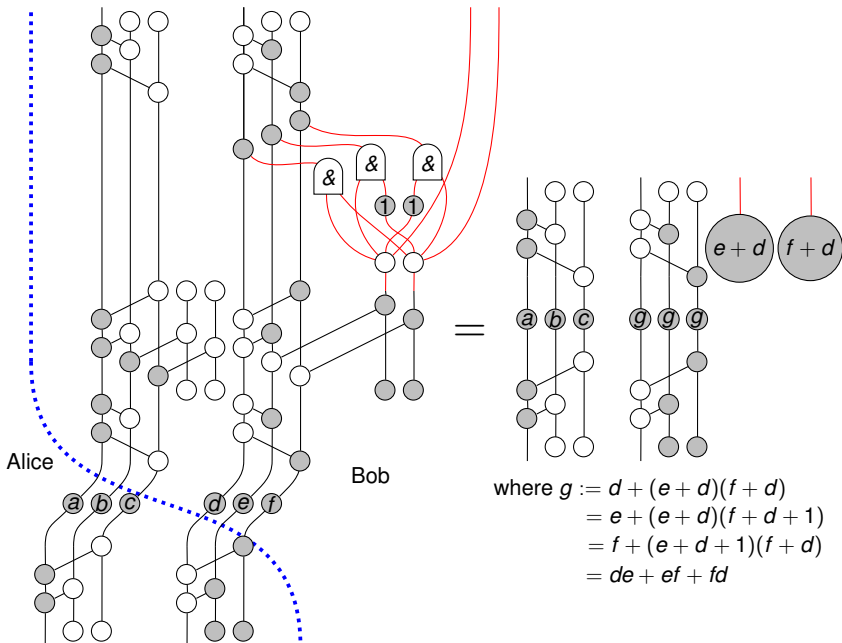
$$X \text{ error } (d, e, f) = (0, 0, 1) \text{ syndrome } (e + d, f + d) = (0, 1)$$

$$X \text{ error } (d, e, f) = (0, 0, 0) \text{ syndrome } (e + d, f + d) = (0, 0)$$

Therefore, we want to apply the correction

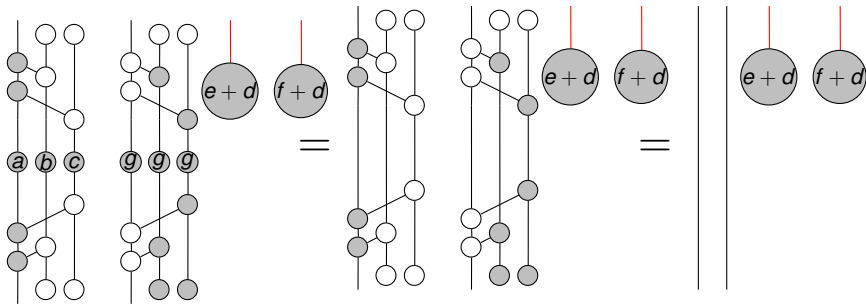
$$\mathbb{F}_2^2 \rightarrow \mathbb{F}_2^{2(3)}; (s, t) \mapsto ((0, 0, 0), (st, s(t + 1)), (s + 1)t)$$

The error correction protocol then has the following form:



If there is at most one X error, then $g = de + ef + fd = 0$.

If furthermore, there are no Z errors:



Problems

The and gate isn't an affine subspace, so it doesn't live in affine coisotropic relations.

Using this formalism we can only *correct* for errors with affine post-processing.

This works for all odd prime qudit stabilizer codes, but only phase-free qubit stabilizer codes (CSS codes)+Z+X gates (no Fourier transform of phase shift)

Future work




Find nice semantics for stabilizer codes with nonlinear classical post-processing.

Undoubled completeness for arbitrary field (following [Booth and Carette, 2022, Poór et al., 2023])





Extending this work to optics/electrical circuits...

Code distance.





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