# Global CNOT Synthesis with Holes 

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## Limitations of NISQ Computers

Circuit

Constraint Topology


## Limitations of NISQ Computers

Circuit

Constraint Topology


## Circuit Synthesis

- Compilation method to overcome NISQ computers shortcommings
- Usually works for circuits made from a specific gate set
- Generates improved circuit, e.g with fewer gates or respecting connectivity constraints, from an efficent representation of the original circuit


## Unitary Decomposition



## Unitary Decomposition

Slicing



## Unitary Decomposition

Slicing


## Unitary Decomposition

Quantum Combs



The problem now becomes how to synthesise a quantum comb

## CNOT Synthesis



## CNOT Combs

Logical Qubits


## CNOT Combs

Temporal Qubits

$\mathcal{H}=\{(1,4),(2,6),(6,7),(4,5)\}$
$p::\{(1,4) \mapsto V,(4,5) \mapsto H,(2,6) \mapsto U,(6,7) \mapsto W\}$

## CNOT Combs



## RowCol - Synthesis Algorithm for CNOT Circuits

## Circuit Representation

$$
\begin{aligned}
& \text { Identity Gate } \\
& |x\rangle-|x\rangle \\
& |y\rangle-|y\rangle \\
& \text { CNOT Gate } \\
& |x\rangle=|x\rangle \\
& |y\rangle-|x \oplus y\rangle
\end{aligned}
$$

$\operatorname{CNOT}(c, t)$ corresponds to $\mathrm{R}(c, t)$
Identity Parity Matrix
$x$
$x^{\prime}$
$y^{\prime}$
1 0

## RowCol - Synthesis Algorithm for CNOT Circuits

## Circuit Representation

CNOT Circuit


CNOT $(c, t)$ corresponds to $\mathrm{R}(c, t)$

Circuit Parity Matrix
$x_{0}^{\prime}$
$x_{0}^{\prime}$
$x_{1}$
$x_{1}^{\prime}$
$x_{2}^{\prime}$
$x_{2}^{\prime}$
$x_{3}^{\prime}$$\left(\begin{array}{cccc}x_{3} \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1\end{array}\right)$

## RowCol - Synthesis Algorithm for CNOT Circuits

Algorithm

- RowCol reduces a parity matrix to the identity by eliminating the row and column for each qubit.
- RowCol can synthesise to constrained architecures:
- Qubit being eliminated has to correspond to non-cutting vertex.
- This is done using Steiner trees.


## RowCol - Synthesis Algorithm for CNOT Circuits

Eliminate Row 1 and Column 1:

$$
\left(\begin{array}{llll}
1 & 0 & 0 & 1 \\
1 & 1 & 1 & 1 \\
0 & 0 & 1 & 1 \\
0 & 0 & 0 & 1
\end{array}\right) \xrightarrow{R_{1}:=R_{0}+R_{1}}\left(\begin{array}{llll}
1 & 0 & 0 & 1 \\
0 & 1 & 1 & 0 \\
0 & 0 & 1 & 1 \\
0 & 0 & 0 & 1
\end{array}\right) \xrightarrow{R_{0}:=R_{3}+R_{0}}\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 1 & 0 \\
0 & 0 & 1 & 1 \\
0 & 0 & 0 & 1
\end{array}\right)
$$

RowCol - Synthesis Algorithm for CNOT Circuits

Eliminate Column 2:

$$
\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 1 & 0 \\
0 & 0 & 1 & 1 \\
0 & 0 & 0 & 1
\end{array}\right) \xrightarrow{R_{1}:=R_{2}+R_{1}}\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 1 \\
0 & 0 & 1 & 1 \\
0 & 0 & 0 & 1
\end{array}\right) \xrightarrow{R_{1}:=R_{3}+R_{1}}\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 1 \\
0 & 0 & 0 & 1
\end{array}\right)
$$

## RowCol - Synthesis Algorithm for CNOT Circuits

Synthesis

Eliminate Column 3:

$$
\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 1 \\
0 & 0 & 0 & 1
\end{array}\right) \xrightarrow{R_{2}:=R_{3}+R_{2}}\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right)
$$

We have reached the identity matrix meaning our synthesis process is over.

## RowCol - Synthesis Algorithm for CNOT Circuits

Row Operations: $R(0,1), R(3,0), R(2,1), R(3,1), R(3,2)$.

New Circuit

Old Circuit


## CombSynth - Synthesis Algorithm for CNOT combs



## CombSynth - Synthesis Algorithm for CNOT combs

$\left.\begin{array}{llllllll}0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 0 \\ 1 \\ 2 \\ 3 \\ 4 \\ 4 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 & 1 & 0 \\ 6 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1\end{array}\right) \quad \mathcal{H}=\{(1,4),(2,6),(6,7),(4,5)\}$

## CombSynth - Synthesis Algorithm for CNOT combs

| 0 |
| :--- |
| 0 |
| 1 |
| 2 |
| 3 |
| 4 |
| 5 |
| 6 |
| 7 |
| 7 |\(\left(\begin{array}{cccccccc}0 \& 1 \& 2 \& 3 \& 4 \& 5 \& 6 \& 7 <br>

0 \& 0 \& 0 \& 1 \& 1 \& 0 \& 1 \& 1 <br>
0 \& 1 \& 0 \& 0 \& 0 \& 0 \& 0 \& 0 <br>
1 \& 0 \& 1 \& 0 \& 0 \& 0 \& 0 \& 0 <br>
1 \& 1 \& 0 \& 1 \& 0 \& 0 \& 0 \& 0 <br>
1 \& 1 \& 0 \& 0 \& 1 \& 0 \& 1 \& 0 <br>
0 \& 0 \& 0 \& 0 \& 0 \& 1 \& 0 \& 0 <br>
1 \& 1 \& 0 \& 1 \& 0 \& 0 \& 1 \& 0 <br>
1 \& 1 \& 0 \& 1 \& 1 \& 0 \& 1 \& 1\end{array}\right) \xrightarrow{ }\)

A temporal qubit can be extracted if its row and column in the full parity matrix can be eliminated by only row operations on the submatrix.

## CombSynth - Synthesis Algorithm for CNOT combs

| Initial sub-matrix | Eliminated sub-matrix | Row operations required for reduction |
| :---: | :---: | :---: |
| $\begin{array}{llllllll}0 & 1 & 2 & 3 & 4 & 5 & 6 & 7\end{array}$ | $\begin{array}{llllllll}0 & 1 & 2 & 3 & 4 & 5 & 6 & 7\end{array}$ |  |
| $0\left(\begin{array}{llllllll}0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0\end{array}\right)$ | $0\left(\begin{array}{llllllll}0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0\end{array}\right)$ |  |
| $7\left(\begin{array}{lllllllll}1 & 1 & 0 & 1 & 1 & 0 & 1 & 1\end{array}\right)$ | $7\left(\begin{array}{lllllllll}1 & 1 & 0 & 1 & 1 & 0 & 1 & 1\end{array}\right)$ | N/A |
| $3\left(\begin{array}{llllllll}1 & 1 & 0 & 1 & 0 & 0 & 0 & 0\end{array}\right)$ | $3\left(\begin{array}{llllllll}1 & 1 & 0 & 1 & 0 & 0 & 0 & 0\end{array}\right)$ |  |
| $\mathcal{H}=\{(1,4),(2,6),(6,7),(4,5)\}$ |  |  |
| $t(0)=0, t(1)=5, t(2)=7, t(3)=3$ |  |  |

## CombSynth - Synthesis Algorithm for CNOT combs



$$
\begin{gathered}
\mathcal{H}=\{(1,4),(2,6),(6,7)\} \\
t(0)=0, t(1)=4, t(2)=7, t(3)=3
\end{gathered}
$$

## CombSynth - Synthesis Algorithm for CNOT combs



$$
\begin{gathered}
\mathcal{H}=\{(1,4),(2,6)\} \\
t(0)=0, t(1)=4, t(2)=6, t(3)=3
\end{gathered}
$$

## CombSynth - Synthesis Algorithm for CNOT combs



$$
\begin{gathered}
\mathcal{H}=\{(1,4)\} \\
t(0)=0, t(1)=2, t(2)=7, t(3)=3
\end{gathered}
$$

## CombSynth - Synthesis Algorithm for CNOT combs



$$
\begin{gathered}
\mathcal{H}=\{(1,4)\} \\
t(0)=0, t(1)=4, t(2)=2, t(3)=3
\end{gathered}
$$

## CombSynth - Synthesis Algorithm for CNOT combs

| Initial sub-matrix |  |  |  |  |  |  |  |  | Eliminated sub-matrix |  |  |  |  |  |  |  |  |  |  | Row operations required for reduction |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 2 | 3 | 4 | 5 | 6 | 7 |  |  | 0 | 1 | 2 | 3 | 4 | 5 | 6 |  |  | $R(0,2), R(1,0)$ |  |
| 0 ( 1 | 1 | 0 | 0 | 0 | 0 | 0 |  |  |  | 1 | 0 | 0 | 0 | 0 | 0 | 0 |  |  |  |  |
| 10 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |  | 0 |  | 0 | 0 | 0 | 0 | 0 |  |  |  |  |
| 21 | 1 | 1 | 0 | 0 | 0 | 0 | 0 |  |  | 0 | 1 | 1 | 0 | 0 | 0 | 0 |  |  |  |  |
|  |  | 0 |  |  |  | 0 |  |  |  |  |  | 0 |  |  |  |  |  |  |  |  |

$$
\begin{gathered}
\mathcal{H}=\{ \} \\
t(0)=0, t(1)=1, t(2)=2, t(3)=3
\end{gathered}
$$

## CombSynth - Synthesis Algorithm for CNOT combs



$$
\begin{gathered}
\mathcal{H}=\{ \} \\
t(0)=0, t(1)=1, t(2)=2, t(3)=3
\end{gathered}
$$

## CombSynth - Synthesis Algorithm for CNOT combs

Row Operations: $R(7,0), R(3,7), R(4,7), R(0,7), R(6,4), R(3,6), R(3,4), R(0,3)$, $R(0,4), R(0,2), R(1,0), R(1,2)$

Holes: $\mathcal{H}=\{(1,4),(2,6),(6,7),(4,5)\}$
Plugs: $p::\{(1,4) \mapsto V,(4,5) \mapsto H,(2,6) \mapsto U,(6,7) \mapsto W\}$

## CombSynth - Synthesis Algorithm for CNOT combs



## Computational Experiments

Slicing



## Computational Experiments

## Experimental Parameters

Architectures<br>■ 9q-square<br>- 16-square<br>■ regetti_16q_aspen<br>- bm_qx5<br>■ ibm_q20_tokyo<br>CNOT Count<br>- 4<br>- 8<br>- 16<br>- 32<br>- 64<br>- 128<br>- 256<br>- 512<br>- 1024

Non-CNOT Proportion

- $5 \%$
- $15 \%$
- $25 \%$
- $50 \%$


## Computational Experiments

Graph Selection





## Computational Experiments

Results

| Architectures | $5 \%$ |  |
| :---: | :---: | :---: |
|  | Non-CNOT Gates |  |
|  | Comb | Slice |
| 9q-square | $-43.1 \%$ | $80.79 \%$ |
| 16q-square | $13.11 \%$ | $344.5 \%$ |
| regetti_ <br> 16q_aspen | $47.31 \%$ | $555.4 \%$ |
| bm_qx5 | $32.27 \%$ | $461.9 \%$ |
| ibm_q20_tokyo | $33.17 \%$ | $393.1 \%$ |

## Computational Experiments

Results

| Architectures | $15 \%$ |  |
| :---: | :---: | :---: |
|  | Non-CNOT Gates |  |
|  | Comb | Slice |
| 9q-square | $34.12 \%$ | $208.7 \%$ |
| 16q-square | $154.4 \%$ | $511.1 \%$ |
| regetti_ <br> $16 q \_a s p e n ~$ | $231.2 \%$ | $893.1 \%$ |
| bm_q×5 | $197.2 \%$ | $698.8 \%$ |
| ibm_q20_tokyo | $183.6 \%$ | $481.0 \%$ |

## Computational Experiments

Results

| Architectures | $25 \%$ |  |
| :---: | :---: | :---: |
|  | Non-CNOT Gates |  |
|  | Comb | Slice |
| 9q-square | $91.93 \%$ | $255.9 \%$ |
| 16q-square | $263.9 \%$ | $564.6 \%$ |
| regetti_ <br> $16 q \_a s p e n ~$ | $379.6 \%$ | $1027 \%$ |
| bm_qx5 | $322.4 \%$ | $783.6 \%$ |
| ibm_q20_tokyo | $289.3 \%$ | $500.2 \%$ |

## Computational Experiments

Results

| Architectures | $50 \%$ |  |
| :---: | :---: | :---: |
|  | Non-CNOT Gates |  |
|  | Comb | Slice |
| 9q-square | $182.1 \%$ | $306.8 \%$ |
| 16q-square | $437.0 \%$ | $606.9 \%$ |
| regetti_ <br> 16q_aspen | $614.9 \%$ | $1119 \%$ |
| bm_qx5 | $527.8 \%$ | $837.7 \%$ |
| ibm_q20_tokyo | $440.7 \%$ | $498.2 \%$ |

## Summary and Conclusion

- Proposed using quantum combs as an alternative to slicing for generalising circuit synthesis
- Introduced CombSynth, a synthesis algorithm for CNOT combs based on RowCol
- Tested CombSynth against RowCol with slicing and found it performed better on a wide variety of experimental parameters.

Future Work

Other Gate Sets



## Future Work

Investigate how synthesis with quantum combs relates to compilation methods that don't use slicing:

- ZX Circuit Extraction: Duncan, Kissinger, Perdrix and van de Wetering 2019
- Lazy Synthesis: Martiel and Goubault de Brugière 2020

